

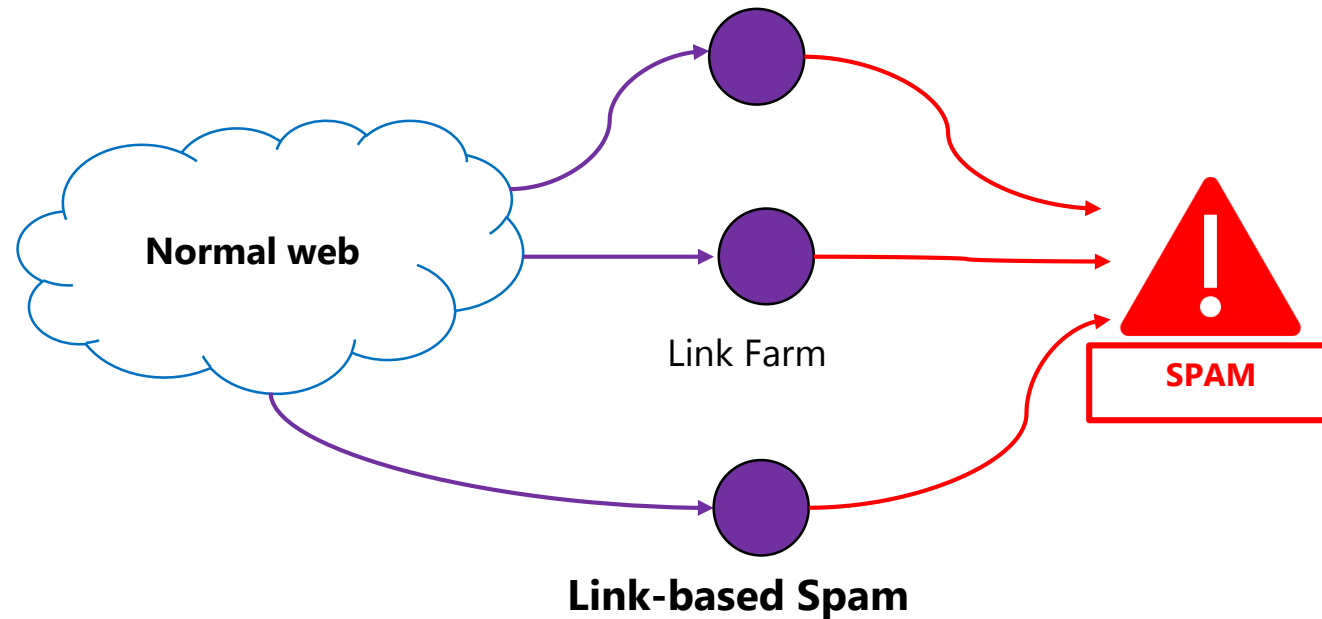


# **Scalable Anti-TrustRank with Qualified Site-level Seeds for Link-based Web Spam Detection**

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# Link-based Web Spam Detection

- **Web Spam**
  - spurious links to get higher-than-deserved rankings.
- Web spam detection algorithms exploit the **hyperlink structure**.



# Contributions

- Collect and share **two real-world web graphs** with labels
- **Two-level analysis** of link spam
  - Page-level graph and site-level graph
  - ATR is useful to detect real-world link spam
- Effective and scalable **site-level seeding methodology** for ATR
- **Asynchronous ATR** significantly reduces the computational cost of ATR

# Real-world Web Graphs

- Crawled by the NAVER search engine (<https://www.naver.com/>)

		page-level graph $G$	site-level graph $H$
W1	No. of normal nodes	797,718 (93.15%)	39,809 (68.63%)
	No. of spam nodes	47,301 (5.52%)	7,954 (13.71%)
	No. of undefined nodes	11,385 (1.33%)	10,239 (17.66%)
	No. of total nodes	856,404	58,002
	No. of labeled edges	3,929,401 (99.33%)	83,351 (85.67%)
	No. of edges	3,955,939	97,294
W2	No. of normal nodes	797,018 (91.20%)	39,984 (67.32%)
	No. of spam nodes	65,259 (7.47%)	8,846 (14.89%)
	No. of undefined nodes	11,684 (1.34%)	10,561 (17.78%)
	No. of total nodes	873,961	59,391
	No. of labeled edges	3,952,584 (99.33%)	84,373 (85.68%)
	No. of total edges	3,979,280	98,478

# Site-level Examination

- A set of **human-labeled seeds**
  - An input of a web spam detection method
- Perform a **site-level examination** followed by refinement of page labels.
- Human experts examine web sites instead of pages.
  - All pages inside a spam site are spam.
  - A normal web site may contain some spam pages
    - Exploit the URL structure to label spam pages

# Two-level Analysis of Link Spam

- Most existing methods focus on either **a page-level graph** or **a site-level graph**, and do not consider both of the graphs.
- We **generalize the structure of link spam** by analyzing the characteristics of link spam on the two different levels of graphs.
  - **Practical solutions for large-scale web spam detection problems**

# Edge Classification

- **Page-level Graph**

- Normal pages tend to point to other normal pages (**TrustRank**)
- Spam pages tend to be referred by other spam pages (**Anti-TrustRank**)

		$ \mathcal{E} $	$E( \mathcal{E} )$	conclusion	$p$ -value
$G$	normal $\rightarrow$ normal	3,639,884	3,500,494	$ \mathcal{E}  > E( \mathcal{E} )$	$7.0 \times 10^{-23}$
	normal $\rightarrow$ spam	2,157	208,725	$ \mathcal{E}  < E( \mathcal{E} )$	$7.9 \times 10^{-28}$
	spam $\rightarrow$ normal	73,049	207,807	$ \mathcal{E}  < E( \mathcal{E} )$	$7.2 \times 10^{-55}$
	spam $\rightarrow$ spam	214,311	12,375	$ \mathcal{E}  > E( \mathcal{E} )$	$9.2 \times 10^{-63}$
$H$	normal $\rightarrow$ normal	56,647	57,840	$ \mathcal{E}  \neq E( \mathcal{E} )$	$2.6 \times 10^{-2}$
	normal $\rightarrow$ spam	17,551	11,771	$ \mathcal{E}  > E( \mathcal{E} )$	$5.6 \times 10^{-13}$
	spam $\rightarrow$ normal	4,394	11,418	$ \mathcal{E}  < E( \mathcal{E} )$	$9.1 \times 10^{-28}$
	spam $\rightarrow$ spam	4,759	2,321	$ \mathcal{E}  > E( \mathcal{E} )$	$9.2 \times 10^{-21}$

# Edge Classification

- **Site-level Graph**

- The number of edges **from normal nodes to spam nodes** is also significant as well as the edges **from spam nodes to spam nodes**.

		$ \mathcal{E} $	$E( \mathcal{E} )$	conclusion	$p$ -value
$G$	normal $\rightarrow$ normal	3,639,884	3,500,494	$ \mathcal{E}  > E( \mathcal{E} )$	$7.0 \times 10^{-23}$
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# Edge Classification

- Consider **an incident node of a between-site edge**
  - (i) The site is **normal** and the page is **normal**
  - (ii) The site is **normal** but the page is **spam**
  - (iii) The site is **spam** and the page is **spam**

# Edge Classification

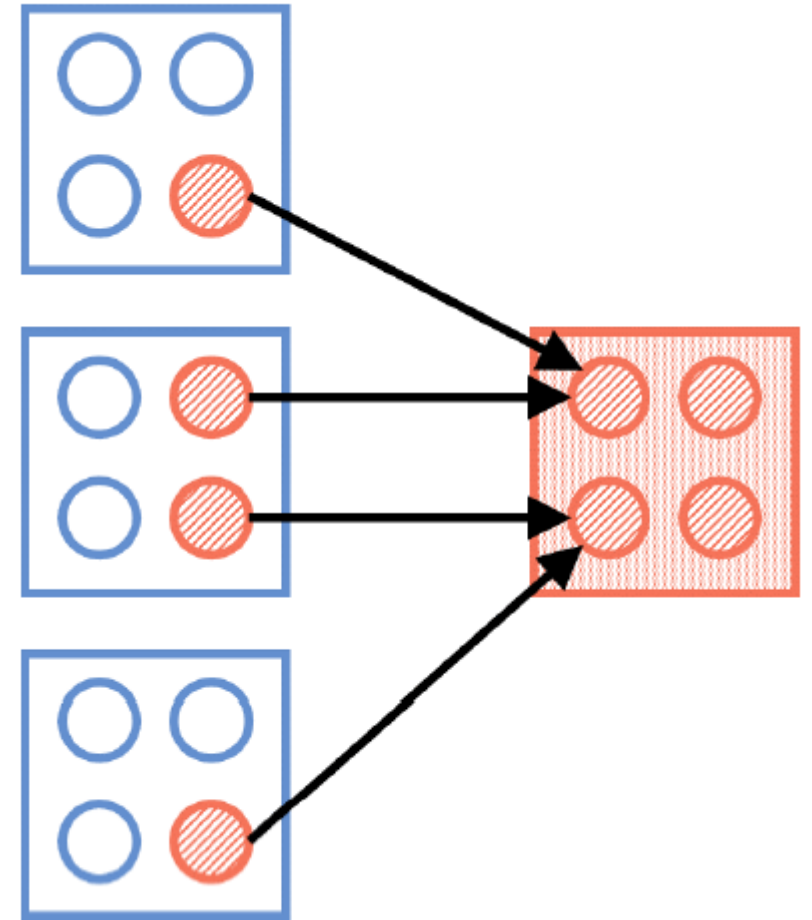
- Three significant edge types: **NSNS, NSSS, SSSS**
  - spam to spam at the page-level graph
- NSSS: normal to spam at the site-level graph

Source		Destination		$ \mathcal{E} $	$E( \mathcal{E} )$	conclusion	$p$ -value
Site	Page	Site	Page				
Normal	Normal	Normal	Normal	857,565	666,284	$ \mathcal{E}  > E( \mathcal{E} )$	$2.0 \times 10^{-20}$
Normal	Normal	Normal	Spam	13	39,750	$ \mathcal{E}  < E( \mathcal{E} )$	$5.5 \times 10^{-17}$
Normal	Normal	Spam	Spam	1,205	5,611	$ \mathcal{E}  < E( \mathcal{E} )$	$5.1 \times 10^{-10}$
Normal	Spam	Normal	Normal	10,825	39,562	$ \mathcal{E}  < E( \mathcal{E} )$	$9.8 \times 10^{-32}$
<b>Normal</b>	<b>Spam</b>	<b>Normal</b>	<b>Spam</b>	52,392	2,357	$ \mathcal{E}  > E( \mathcal{E} )$	$4.9 \times 10^{-55}$
<b>Normal</b>	<b>Spam</b>	<b>Spam</b>	<b>Spam</b>	121,397	336	$ \mathcal{E}  > E( \mathcal{E} )$	$1.7 \times 10^{-85}$
Spam	Spam	Normal	Normal	5,953	7,361	$ \mathcal{E}  < E( \mathcal{E} )$	$1.3 \times 10^{-5}$
Spam	Spam	Normal	Spam	340	453	$ \mathcal{E}  < E( \mathcal{E} )$	$2.6 \times 10^{-3}$
<b>Spam</b>	<b>Spam</b>	<b>Spam</b>	<b>Spam</b>	3,768	67	$ \mathcal{E}  > E( \mathcal{E} )$	$2.0 \times 10^{-52}$

# Web Spam via Two-level Edge Classification

## • Overpost

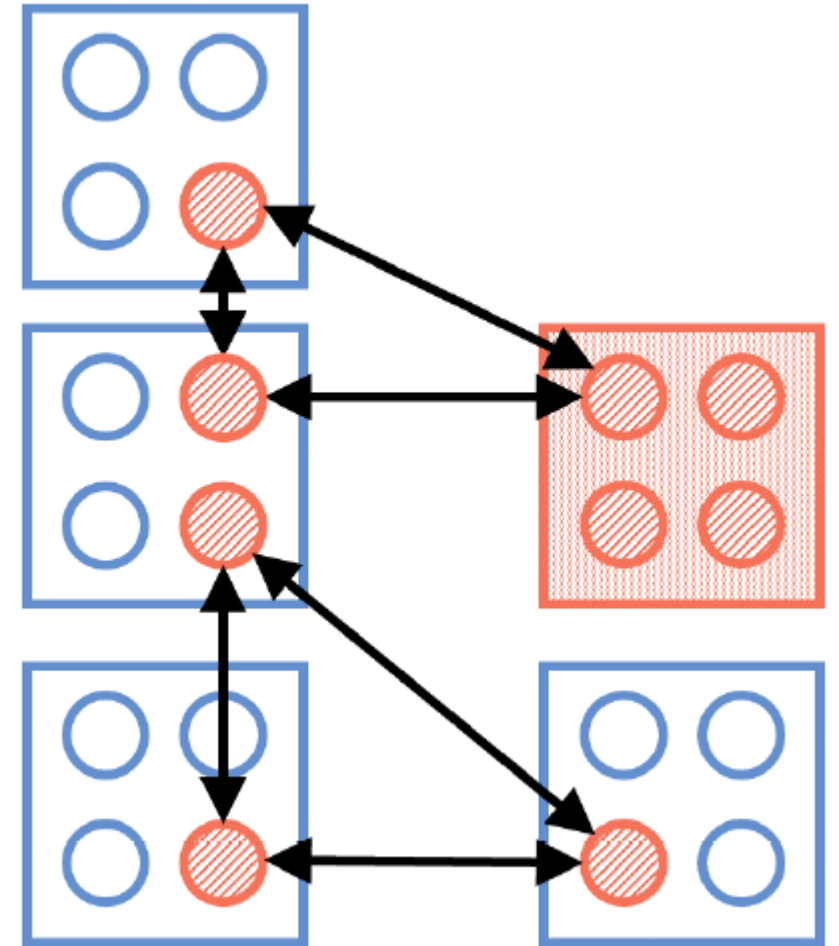
- A spammer makes a lot of postings in different normal sites to intrigue transactions into the targeting spam site.
- The postings are spam pages which contain the links to the spam pages in the spam site.
- This configuration makes the **NSSS** edge type.



# Web Spam via Two-level Edge Classification

## • Hacking

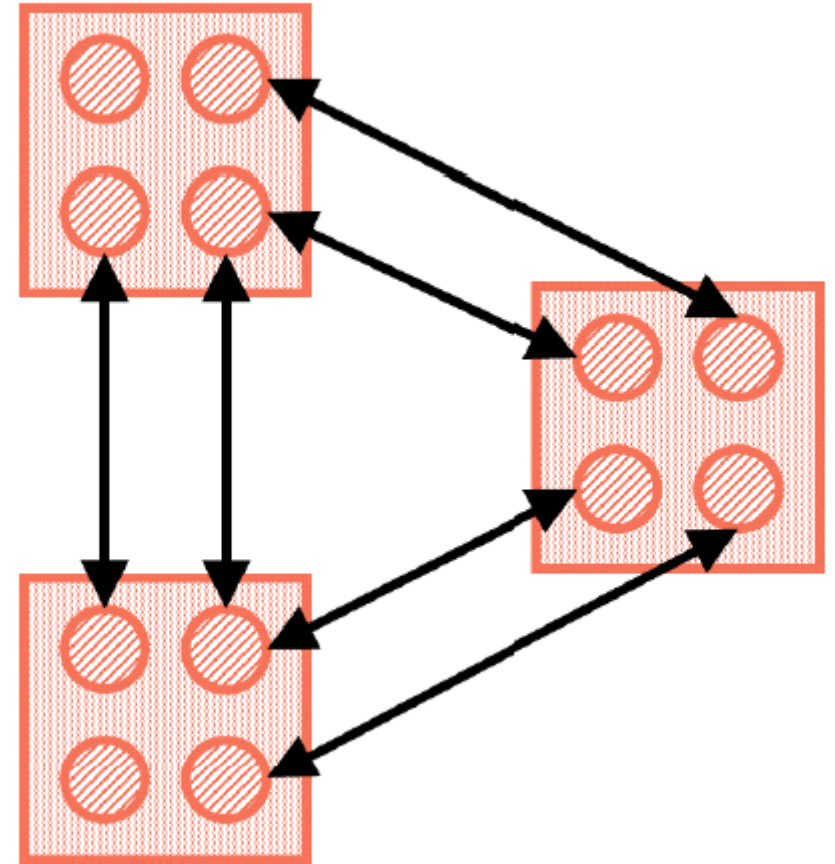
- A spammer hacks normal sites. The spammer makes spam pages in normal sites and the spam pages are linked to other spam pages.
- We can observe the **NSSS** and **NSNS** edges.



# Web Spam via Two-level Edge Classification

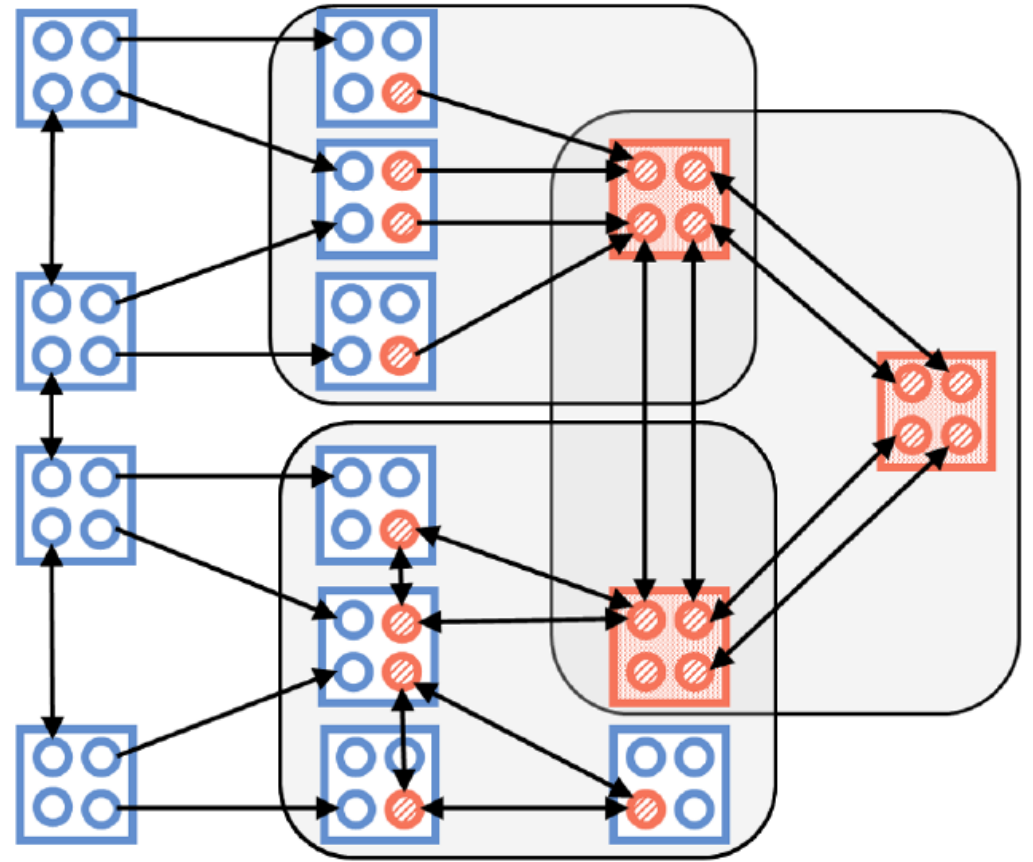
## • Link Farm

- Some spam sites and spam pages are designed to be densely connected with each other to raise PageRank scores so that they can be indexed by a search engine.
- We observe **SSSS** edge types.



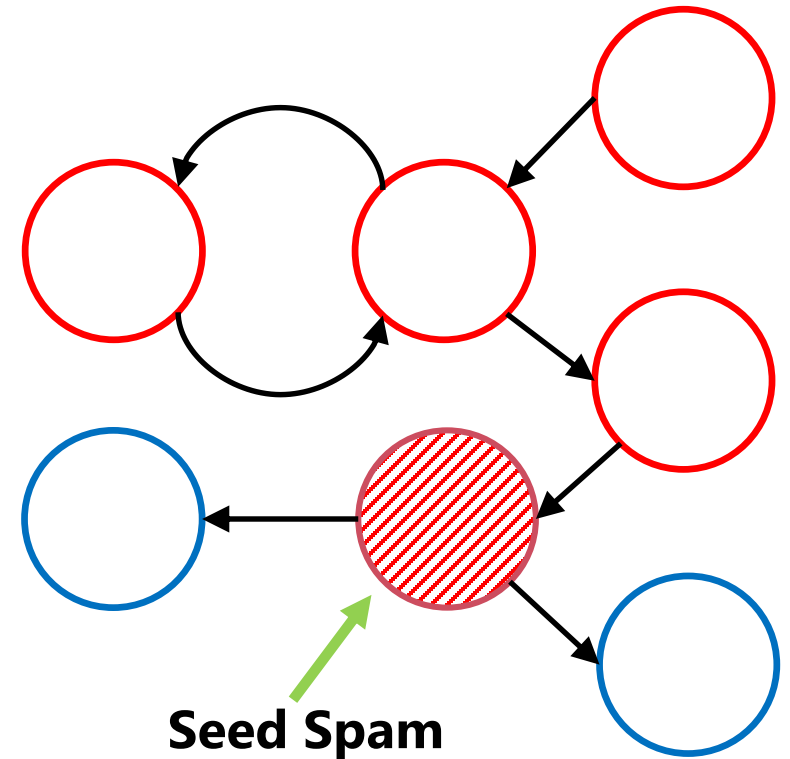
# Web Spam via Two-level Edge Classification

- Real-world link spam can be explained by a combination of the aforementioned **building blocks**.



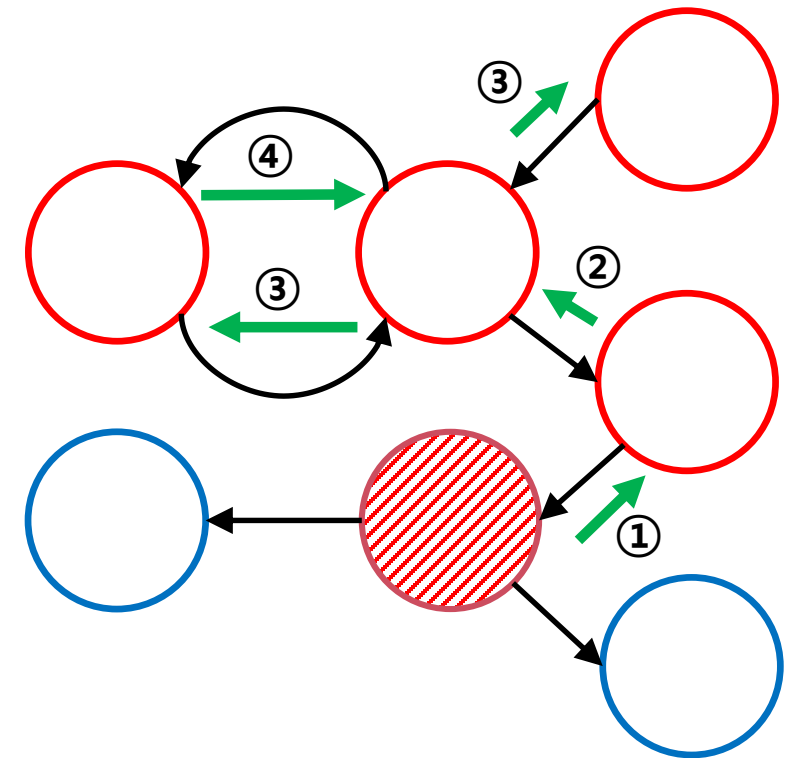
# Anti-TrustRank with Qualified Site-level Seeds

- Anti-TrustRank (ATR)
  - Spam pages are likely to be referenced by other spam pages.
  - **Carefully select seed spam pages.**
  - Assign ATR scores to the seed spam pages.



# Anti-TrustRank with Qualified Site-level Seeds

- Anti-TrustRank (ATR)
  - From the seeds, the ATR scores are propagated to incoming neighbors of the nodes so that **the pages having links to the spam pages end up with having high ATR scores.**
  - Pages with high ATR scores are considered as spam pages.





# Anti-TrustRank with Qualified Site-level Seeds

- The spam seeds should be examined by human experts to get labels.
- Human experts conduct a site-level examination.
- **Represent each site as a feature vector** and build a classifier that predicts the probability of being spam.
- We **prioritize the websites** according to the probability for the **site-level examination**.

# Anti-TrustRank with Qualified Site-level Seeds

- **Our features to model a site**

- **entro-in-p**: the entropy of the indegrees of pages within a site

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*in-p*: indegree of each page in the site  $h$

*out-p*: outdegree of each page in the site  $h$

*dist*: the distances from the site  $h$  to all other reachable sites on  $\bar{H}$

---

**entro-in-p**: entropy of *in-p*

**entro-out-p**: entropy of *out-p*

**mean-dist**: mean of *dist*

**std-dist**: standard deviation of *dist*

**max-dist**: maximum of *dist*

**within-site**: no. of within-site edges

**in-h**: indegree of the site  $h$  on  $H$

**out-h**: outdegree of the site  $h$  on  $H$

**reachability**: no. of reachable sites on  $\bar{H}$

**cluster**: whether  $h$  belongs to a spam cluster

**dmnt-ratio**: max. weight/degree of  $h$  on  $\bar{H}_w$

**no-page**: no. of pages in the site  $h$

**in-page**: no. of pages having an edge to  $h$

**out-page**: no. of pages having an edge from  $h$

**one-hop**: no. of one-hop distant sites on  $\bar{H}$

**two-hop**: no. of two-hop distant sites on  $\bar{H}$

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# Anti-TrustRank with Qualified Site-level Seeds

- **Classification performance of the features**
  - Our features show better performance than node2vec features.

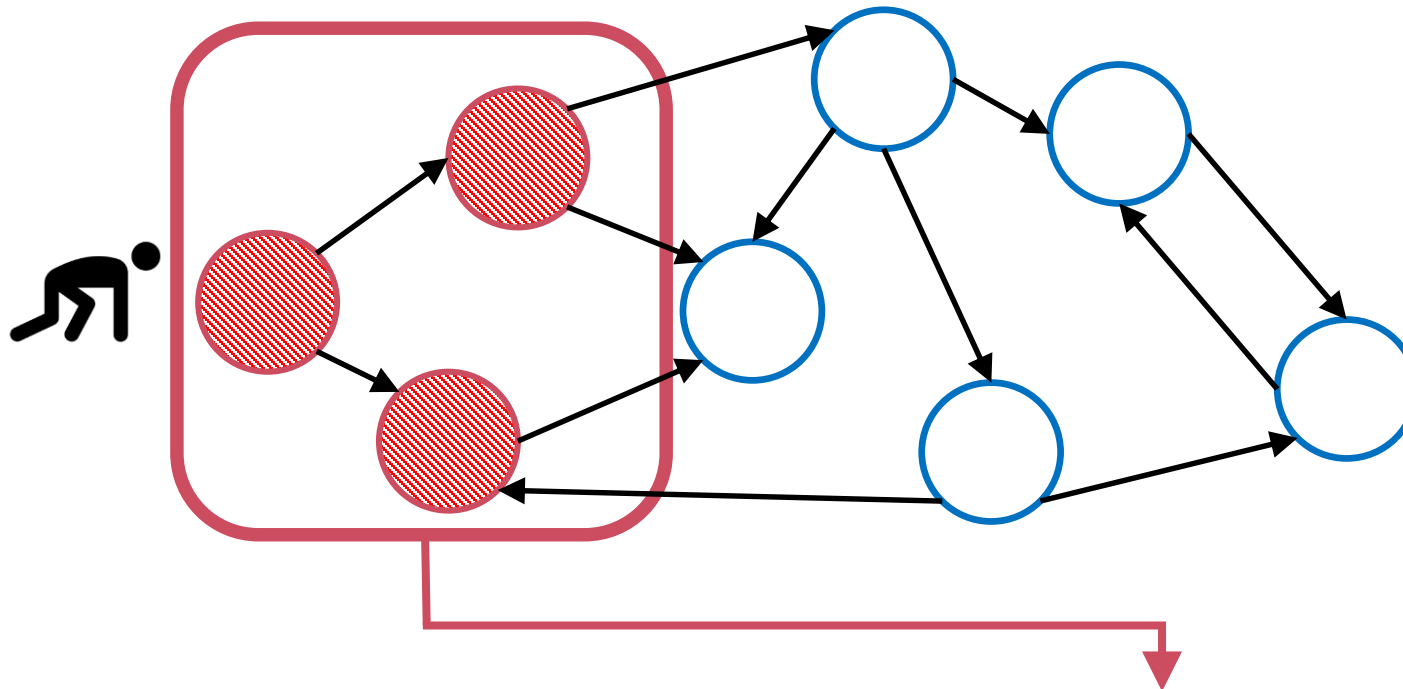
	W1		W2	
	node2vec	Our Features	node2vec	Our Features
Accuracy	83.9%	<b>88.0%</b>	82.7%	<b>88.1%</b>
Normal F1	90.6%	<b>92.1%</b>	89.7%	<b>92.2%</b>
Spam F1	46.1%	<b>86.1%</b>	45.1%	<b>86.1%</b>
Avg. Precision	70.5%	<b>88.8%</b>	70.2%	<b>89.0%</b>
Avg. Recall	66.8%	<b>89.4%</b>	65.7%	<b>89.3%</b>
Avg. F1	68.3%	<b>89.1%</b>	67.4%	<b>89.1%</b>

# Work-Efficient Anti-TrustRank

- Computing **Anti-TrustRank (ATR)** scores is identical to computing the **personalized PageRank (PPR)** scores on the reverse graph.
  - Spam seeds in ATR → personalization set (predefined nodes) in PPR
- We propose **asynchronous Anti-TrustRank algorithms**
  - Reduce the computational cost of the traditional ATR algorithm
  - Without degrading performance in spam detection
  - Convergence analysis

# Personalized PageRank

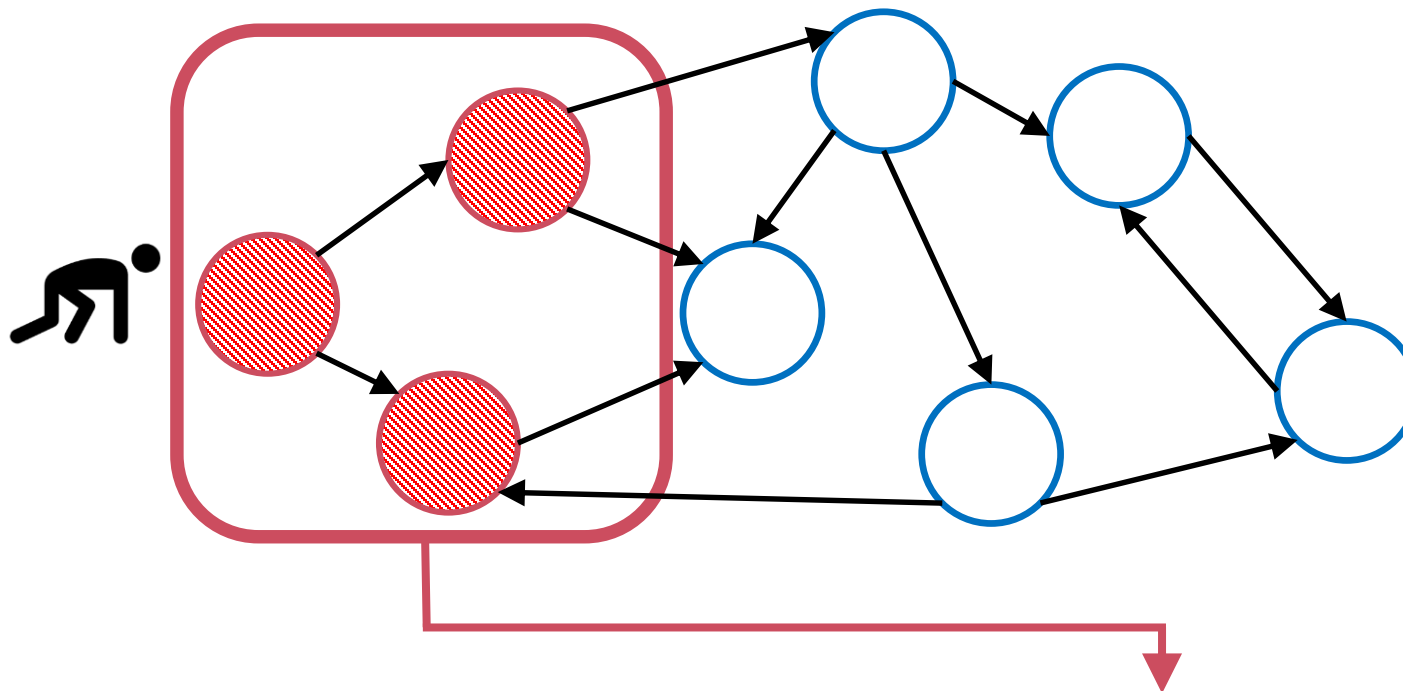
- Randomly jump to one of **the predefined nodes**.



$$x = \alpha P^T x + (1 - \alpha) e_s \quad (e_s : \text{Personalized vector})$$

# Anti-TrustRank

- Randomly jump to one of **spam seeds** on the reverse graph.

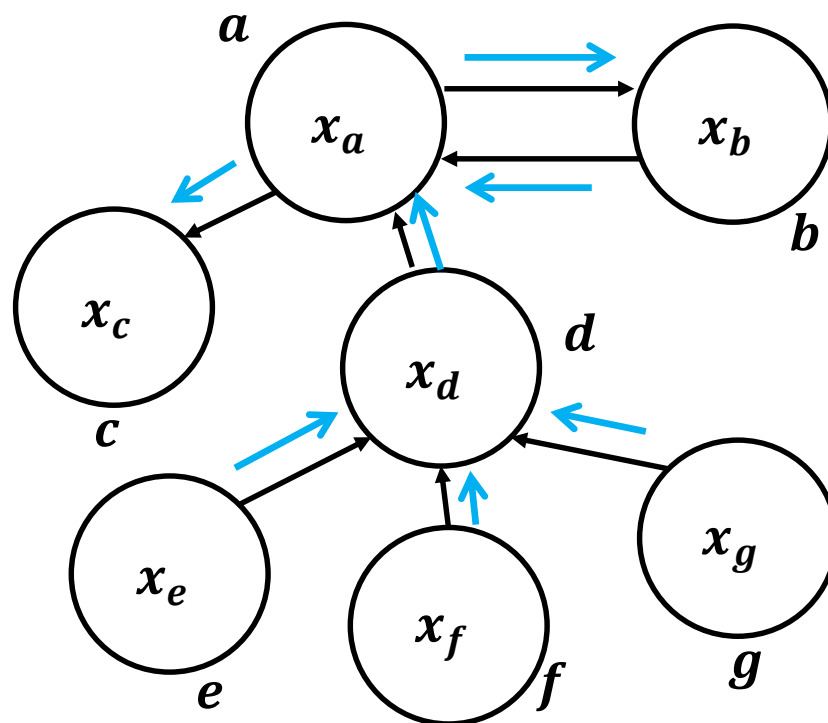


$$x = \alpha P^T x + (1 - \alpha) e_s \quad (e_s : \text{Personalized vector})$$

# Synchronous Anti-TrustRank

## SYNC ATR

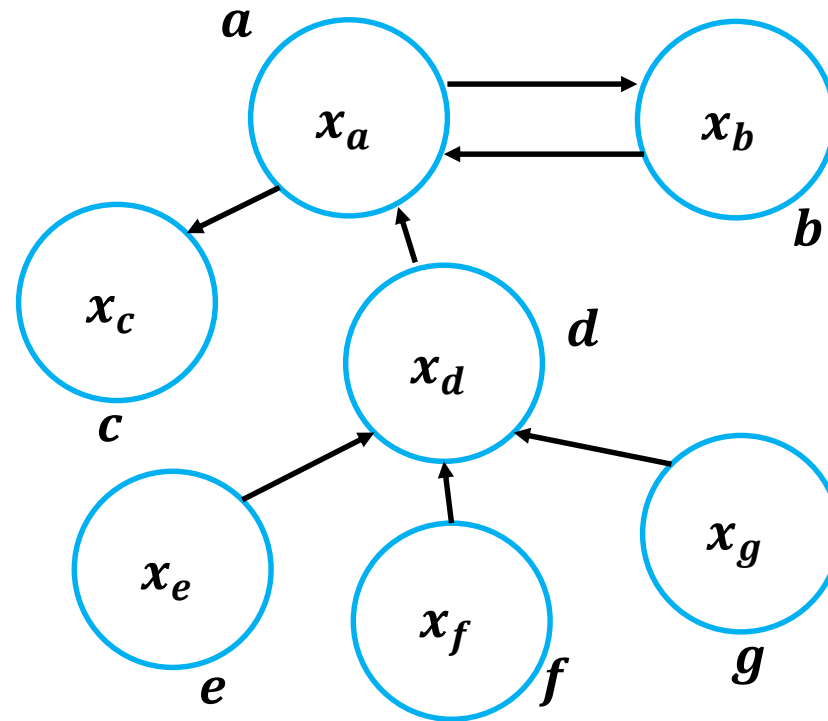
The ATR scores are updated after all the nodes re-compute the ATR scores.



# Synchronous Anti-TrustRank

## SYNC ATR

The ATR scores are updated after all the nodes re-compute the ATR scores.

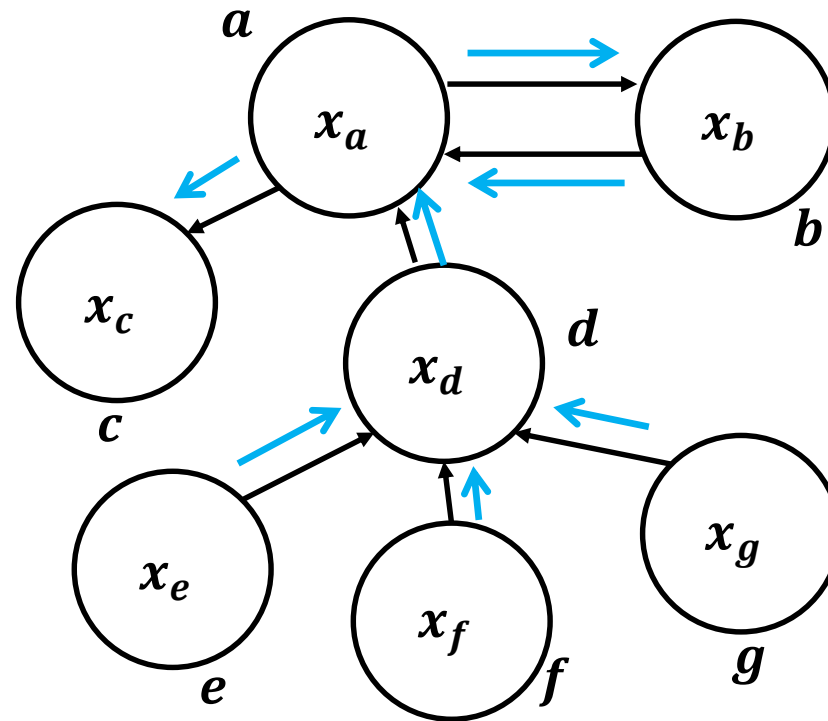




# Synchronous Anti-TrustRank

## SYNC ATR

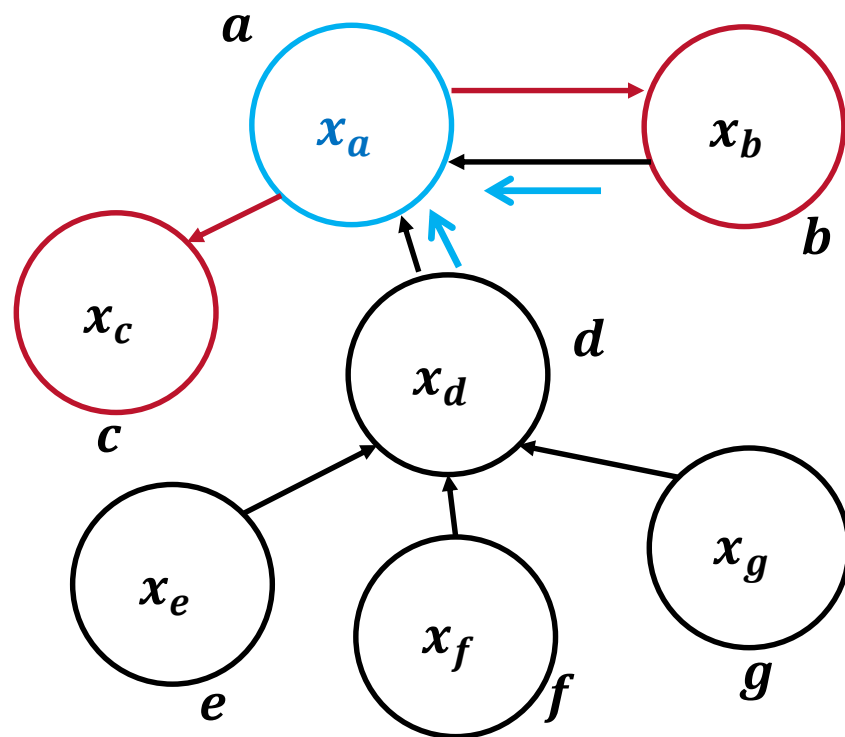
The ATR scores are updated after all the nodes re-compute the ATR scores.



# Asynchronous Anti-TrustRank

## ASYNC ATR

**Worklist** A set of nodes whose ATR scores need to be updated



**Worklist**

[b, c, d, e, f, g, **b**, **c**]

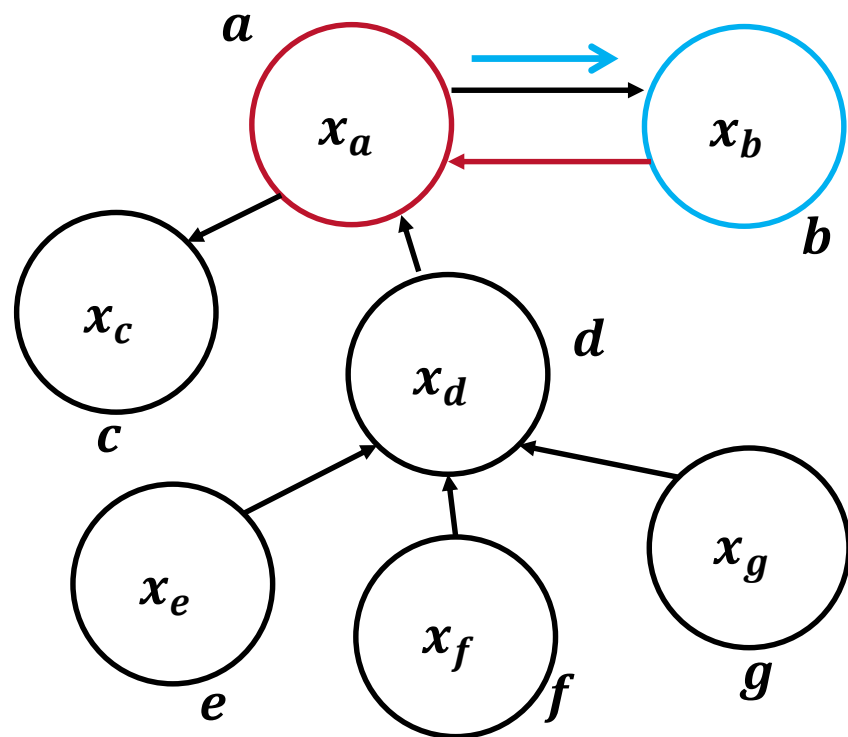
**Pop a**

**Push b, c**

# Asynchronous Anti-TrustRank

## ASYNC ATR

**Worklist** A set of nodes whose ATR scores need to be updated



**Worklist**

[c, d, e, f, g, **b**, **c**, **a**]

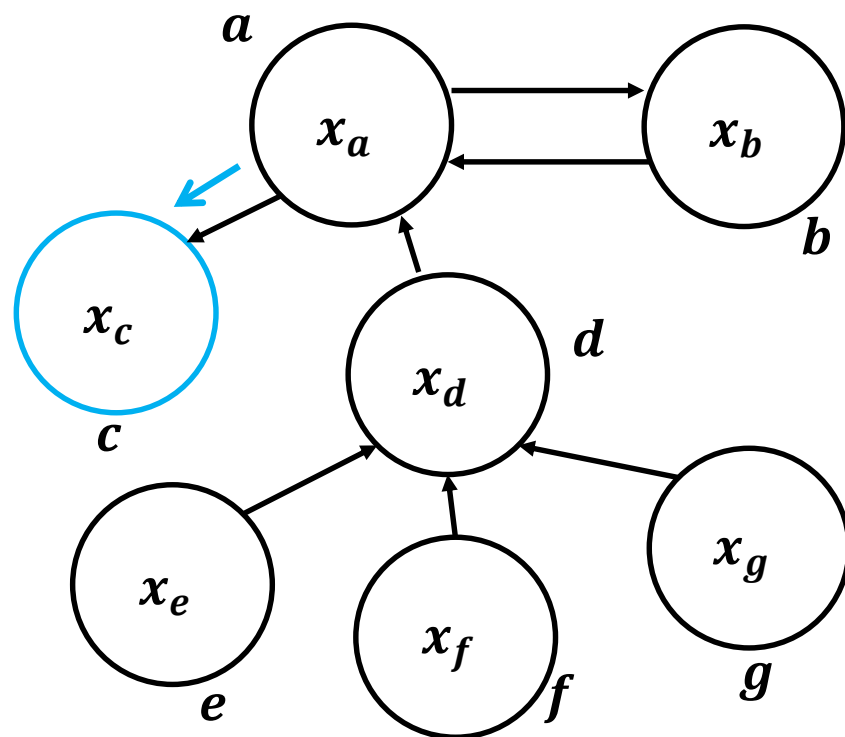
**Pop b**

**Push a**

# Asynchronous Anti-TrustRank

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**Worklist**

[d, e, f, g, **b**, **c**, **a**]

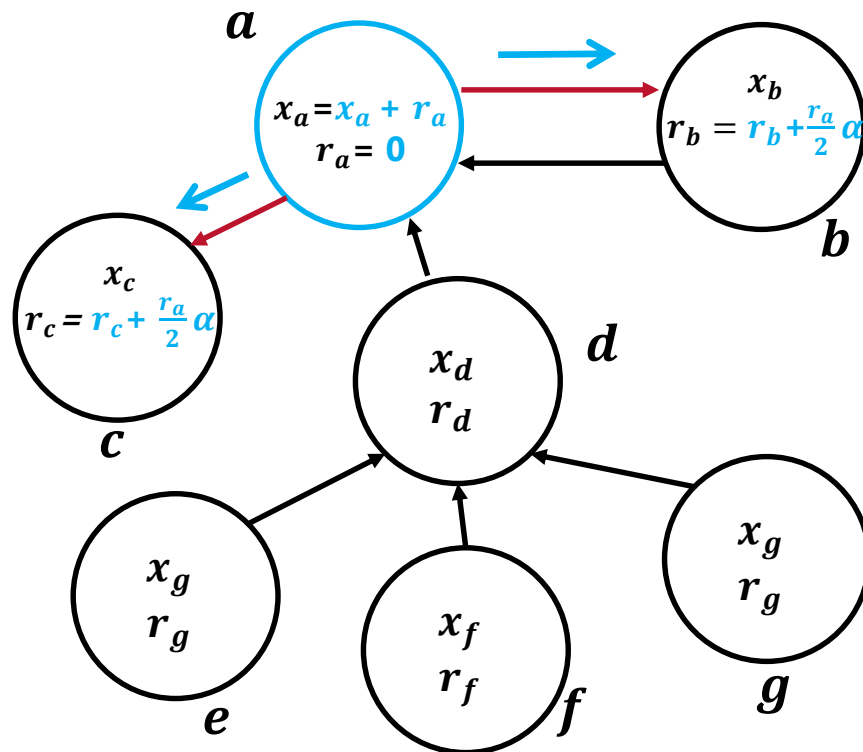
**Pop c**

# Residual-based Asynchronous Anti-TrustRank

## RASYNC ATR

**new ATR = current ATR + current residual** (explicitly maintain the residual of each node)

**Filtering out unnecessary work** in the worklist



**Worklist**

[b, c, d, e, f, g]

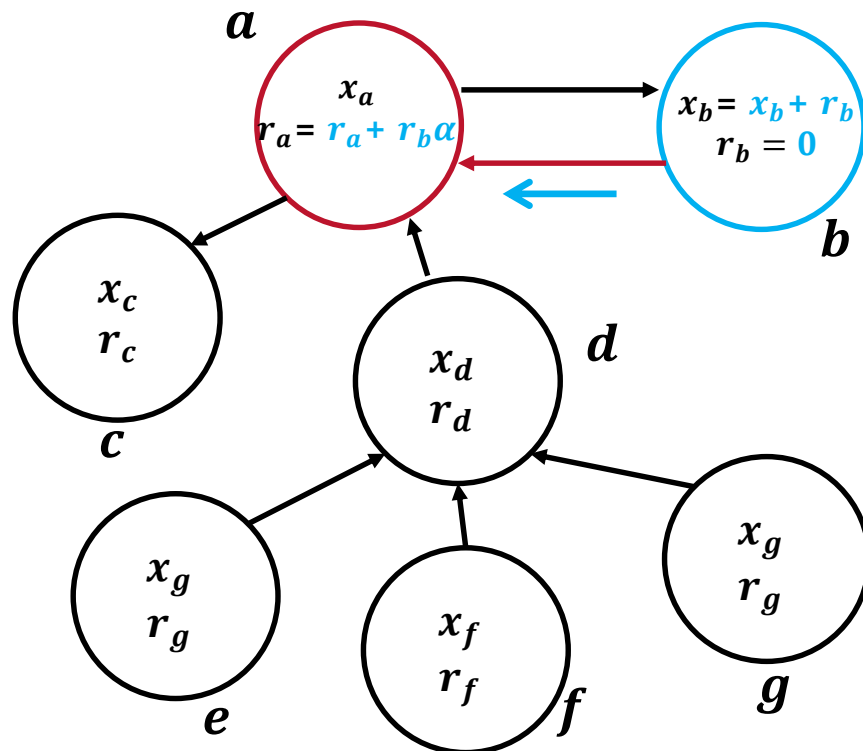
**Pop a**

# Residual-based Asynchronous Anti-TrustRank

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**Worklist**

[c, d, e, f, g, **a**]

**Pop b**

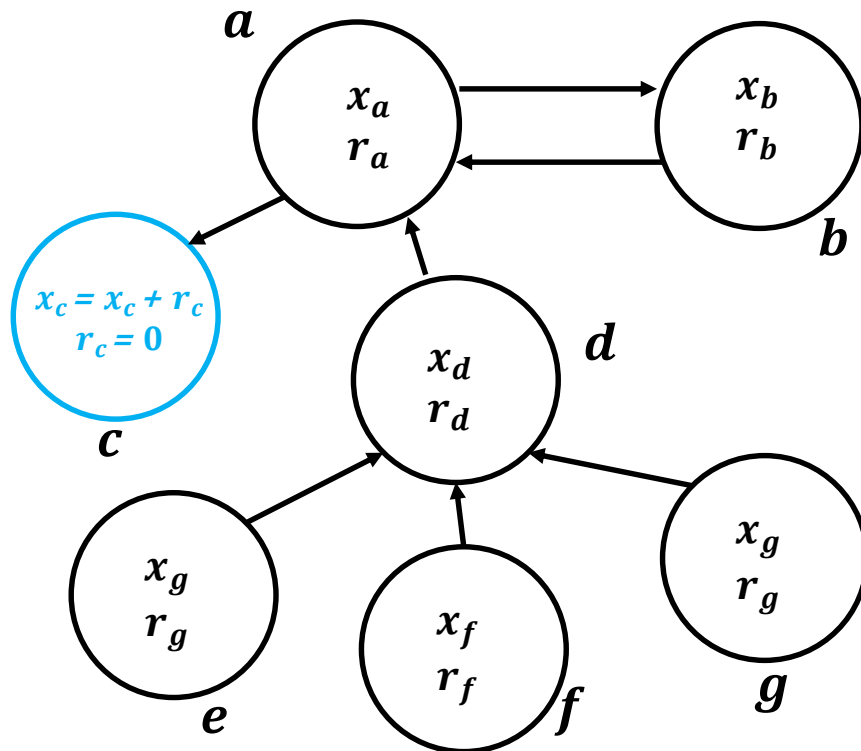
**Push a**

# Residual-based Asynchronous Anti-TrustRank

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**Filtering out unnecessary work** in the worklist



**Worklist**

[d, e, f, g, **a**]

**Pop c**

# Convergence of Asynchronous Anti-TrustRank

The Anti-TrustRank  $\mathbf{x}$  is computed as follows:

$$\mathbf{x} = \alpha \mathbf{P}^T \mathbf{x} + (1 - \alpha) \mathbf{e}_s.$$

Where  $\mathbf{P}$  is a row-stochastic matrix ( $\mathbf{P} = \mathbf{D}^{-1} \mathbf{A}$ ) and  $\mathbf{e}_s$  is the personalized vector. This is the linear system:

$$(\mathbf{1} - \alpha \mathbf{P}^T) \mathbf{x} = (1 - \alpha) \mathbf{e}_s.$$

and the residual:

$$\mathbf{r} = (1 - \alpha) \mathbf{e}_s - (\mathbf{1} - \alpha \mathbf{P}^T) \mathbf{x} = \alpha \mathbf{P}^T \mathbf{x} + (1 - \alpha) \mathbf{e}_s - \mathbf{x}.$$

When the  $j$ -th node is processed, the residual is decreased by  $r_j^k (1 - \alpha)$ .

$$\mathbf{e}^T \mathbf{r}^{(k+1)} = \mathbf{e}^T \mathbf{r}^{(k)} - r_j^k (1 - \alpha).$$



# Asynchronous Anti-TrustRank Algorithms

## Asynchronous Anti-TrustRank

- Require much fewer Anti-TrustRank updates as well as arithmetic operations with the same precision by maintaining a working set.

## Residual-based Asynchronous Anti-TrustRank

- Significantly reduces the number of arithmetic computations.
- Able to effectively reduce the size of the working set by filtering out unnecessary computations.

# Experimental Results

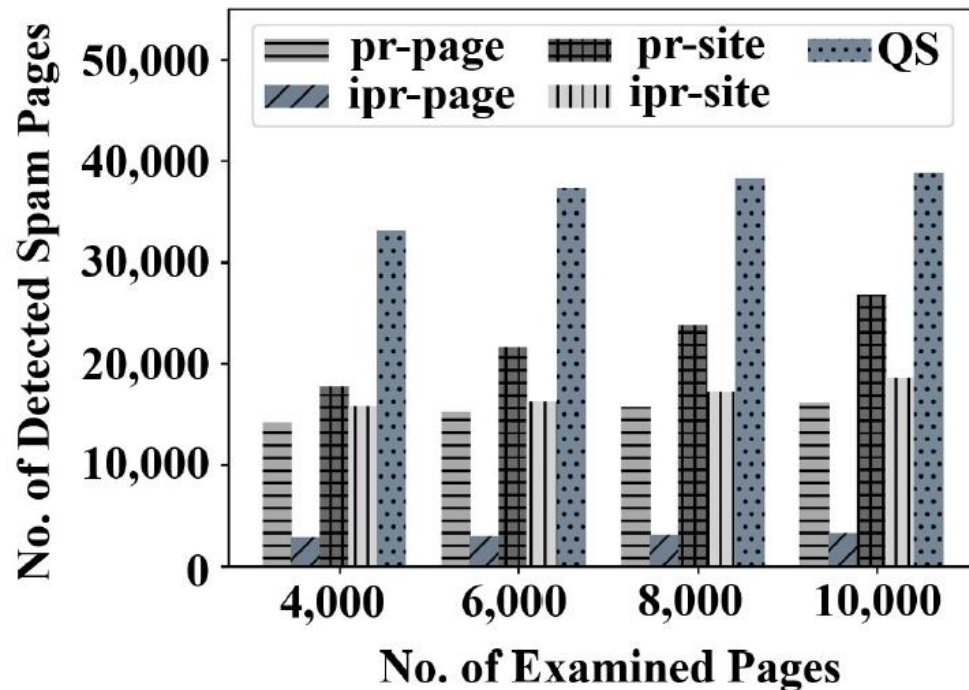
- Performance in web spam detection

- Our method (**QS**) significantly outperforms other methods.

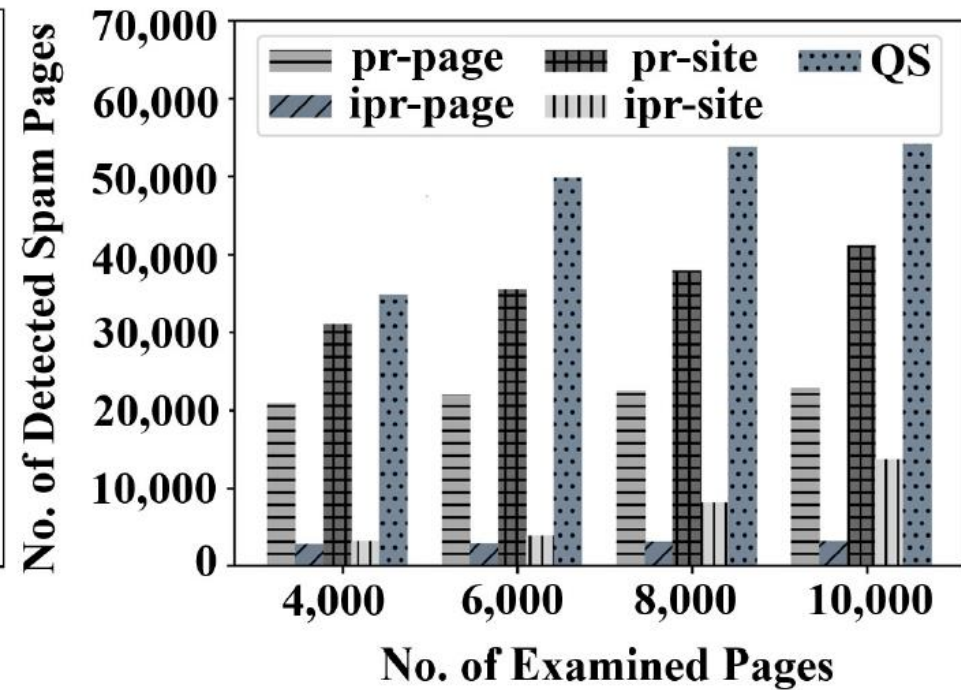
No. of Examined Pages		<i>lfeat</i>	<i>nvec</i>	<i>trust</i>	<i>pr-page</i>	<i>ipr-page</i>	<i>pr-site</i>	<i>ipr-site</i>	<b>QS</b>
4,000 (0.47% examined)	Accuracy	60.80%	94.50%	26.33%	96.00%	94.73%	96.41%	96.25%	<b>98.22%</b>
	F1 score	15.90%	5.80%	13.04%	45.67%	11.22%	53.90%	49.95%	<b>81.52%</b>
	Precision	9.00%	68.30%	6.98%	95.56%	98.43%	96.14%	99.05%	97.67%
	Recall	66.20%	3.00%	99.83%	30.01%	5.95%	37.45%	33.39%	69.95%
6,000 (0.70% examined)	Accuracy	89.20%	94.60%	27.39%	96.12%	94.75%	96.86%	96.31%	<b>98.71%</b>
	F1 score	21.40%	22.10%	13.21%	48.20%	11.75%	61.98%	51.03%	<b>87.22%</b>
	Precision	18.00%	57.00%	7.07%	95.51%	98.27%	96.34%	99.01%	97.60%
	Recall	26.40%	13.70%	99.87%	32.23%	6.25%	45.69%	34.37%	78.83%
10,000 (1.17% examined)	Accuracy	84.30%	94.40%	35.02%	96.21%	94.78%	97.47%	96.58%	<b>98.88%</b>
	F1 score	21.70%	30.90%	14.53%	50.16%	12.77%	71.46%	56.28%	<b>89.12%</b>
	Precision	15.10%	49.40%	7.84%	95.14%	98.09%	96.75%	98.89%	97.42%
	Recall	38.80%	22.50%	99.83%	34.06%	6.83%	56.65%	39.33%	82.13%

# Experimental Results

- No. of detected spam pages of the ATR algorithm with different seeding methods
  - Our seeding method (**QS**) detects the largest number of spam pages.



(a) W1 dataset



(b) W2 dataset

# Experimental Results

- **async** and **rasync** save much computation compared to **sync**.
- **rasync** reduces the number of arithmetic operations compared to **async**.

		<i>sync</i>	<i>async</i>	<i>rasync</i>
$\epsilon = 10^{-4}$	No. of Detected Spam Pages	33,088	33,029	33,029
	F1 Score	81.52 %	81.67 %	81.67 %
	No. of ATR updates	51,384,240	46,680	46,454
$e = 4000$	No. of Arithmetics	578,549,460	11,170,087	1,765,129
	Run Time (milliseconds)	7,596	339	87
<hr/>				
$\epsilon = 10^{-8}$	No. of Detected Spam Pages	33,088	33,088	33,088
	F1 Score	81.52 %	81.52 %	81.52 %
	No. of ATR updates	100,199,268	83,961	83,972
$e = 4000$	No. of Arithmetics	1,128,171,447	13,009,448	2,673,169
	Run Time (milliseconds)	14,952	358	99

# Experimental Results

- Run Time (milliseconds) of the algorithms
  - **rasync** is the fastest method.

		<i>sync</i>	<i>async</i>	<i>rasync</i>	<i>bstab</i>	<i>brppr</i>
W1	$e=4,000, \epsilon=10^{-4}$	7,596	339	87	566	678
	$e=4,000, \epsilon=10^{-8}$	14,952	358	99	1,217	680
	$e=10,000, \epsilon=10^{-4}$	7,678	350	98	678	822
	$e=10,000, \epsilon=10^{-8}$	14,628	374	111	1,775	829
W2	$e=4000, \epsilon=10^{-4}$	6,526	556	148	821	726
	$e=4,000, \epsilon=10^{-8}$	13,841	1,205	374	1,926	742
	$e=10,000, \epsilon=10^{-4}$	6,212	607	169	707	968
	$e=10,000, \epsilon=10^{-8}$	13,174	1,406	453	1,546	948

# Experimental Results

- Parallel sync, async, and rasync on distributed machines
  - **rasync** is the fastest method.

Data Information			Run Time (minutes)		
No. of nodes	No. of edges	Size of $\mathcal{S}$	<i>sync</i>	<i>async</i>	<i>rasync</i>
59,180,800	82,824,237	2,340,940	86	94	37
152,595,632	274,392,463	3,329,026	191	162	69
57,135,532	732,008,321	4,381,555	516	351	121
556,047,762	1,207,335,482	5,016,499	>2,116	>1,413	163

# Conclusion

- We develop a **site-level seeding methodology** for the ATR algorithm, which leads to remarkably boosting up the performance of the ATR algorithm.
- We design a work-efficient **asynchronous ATR algorithm** which significantly reduces the computational cost of the traditional ATR method while guaranteeing convergence.
- Our methodologies can be **integrated into other spam detection models** in practice, e.g., considering both TrustRank and Anti-TrustRank.

# Big Data Lab

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